

Folds in Haskell

Mark P Jones

Portland State University

Folds!

- ◆ A list xs can be built by applying the $(:)$ and $[]$ operators to a sequence of values:

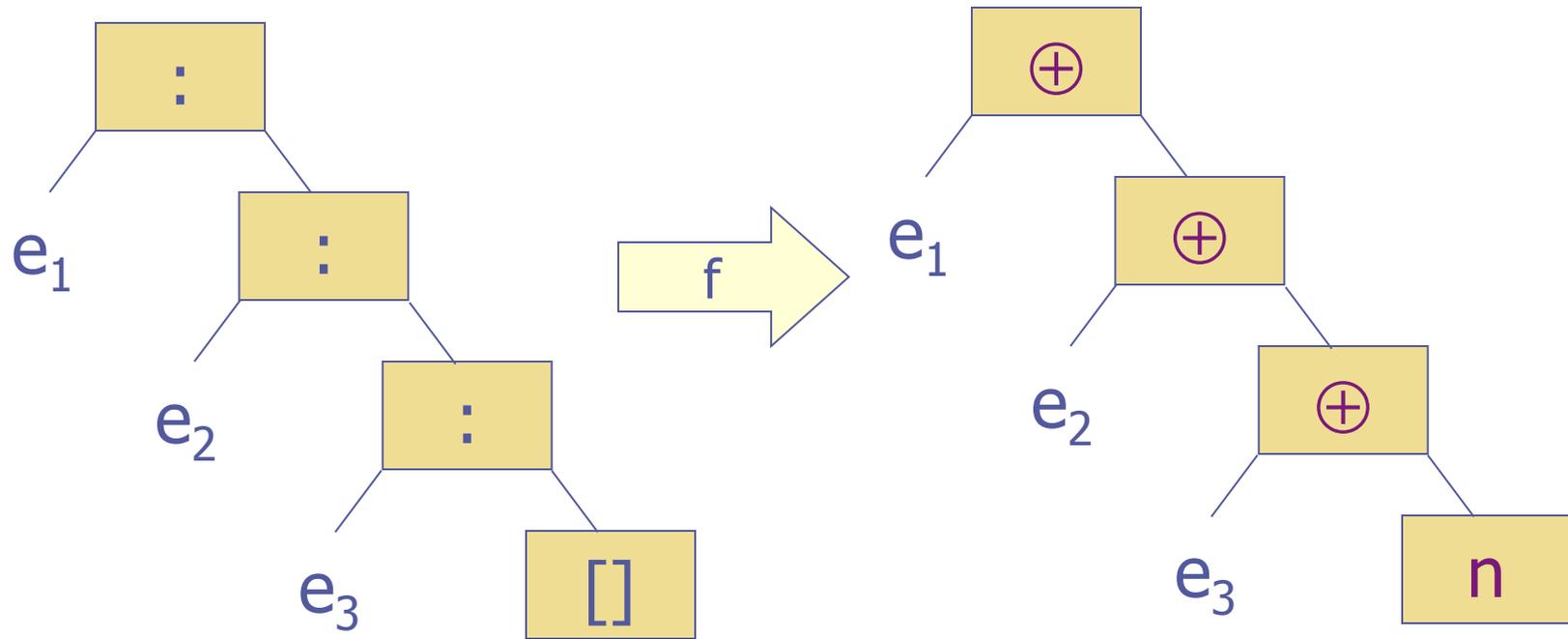
$$xs = x_1 : x_2 : x_3 : x_4 : \dots : x_k : []$$

- ◆ Suppose that we are able to replace every use of $(:)$ with a binary operator (\oplus) , and the final $[]$ with a value n :

$$xs = x_1 \oplus x_2 \oplus x_3 \oplus x_4 \oplus \dots \oplus x_k \oplus n$$

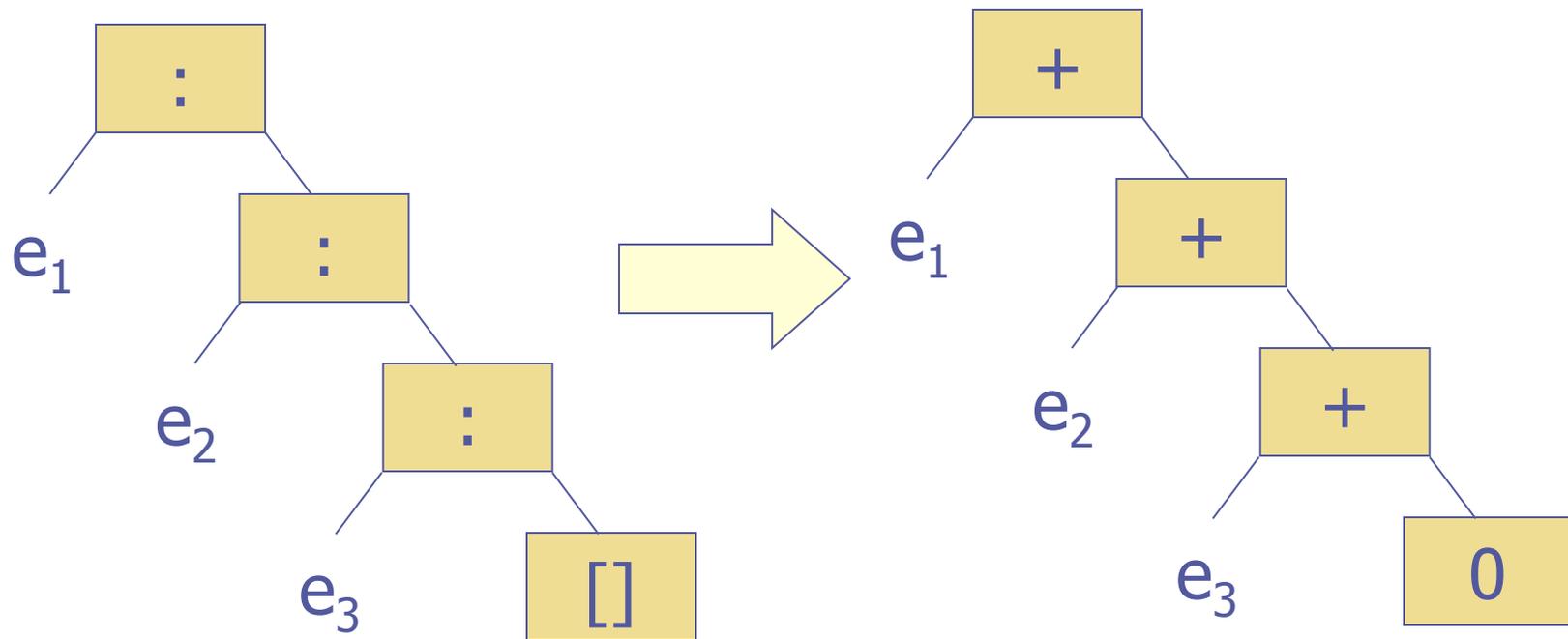
- ◆ The resulting value is called $\text{fold } (\oplus) n xs$
- ◆ Many useful functions on lists can be described in this way.

Graphically:



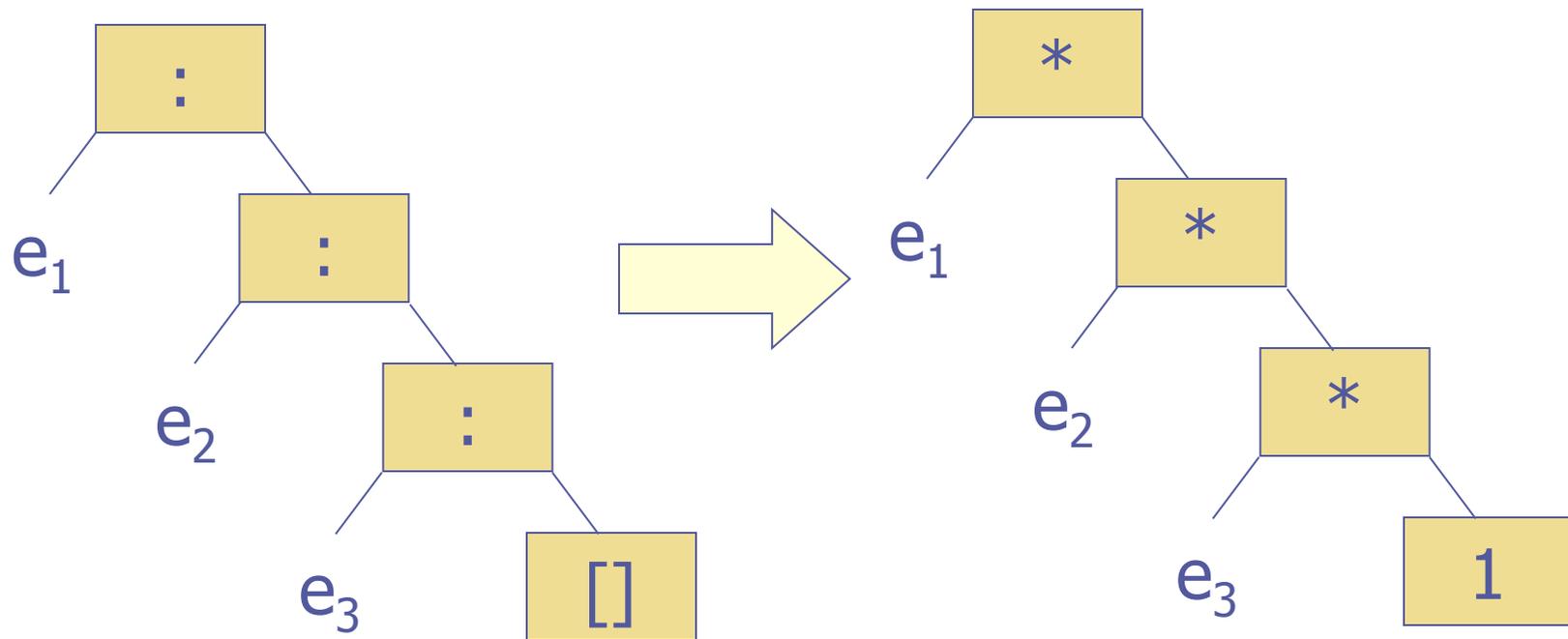
$$f = \text{foldr } (\oplus) n$$

Example: sum



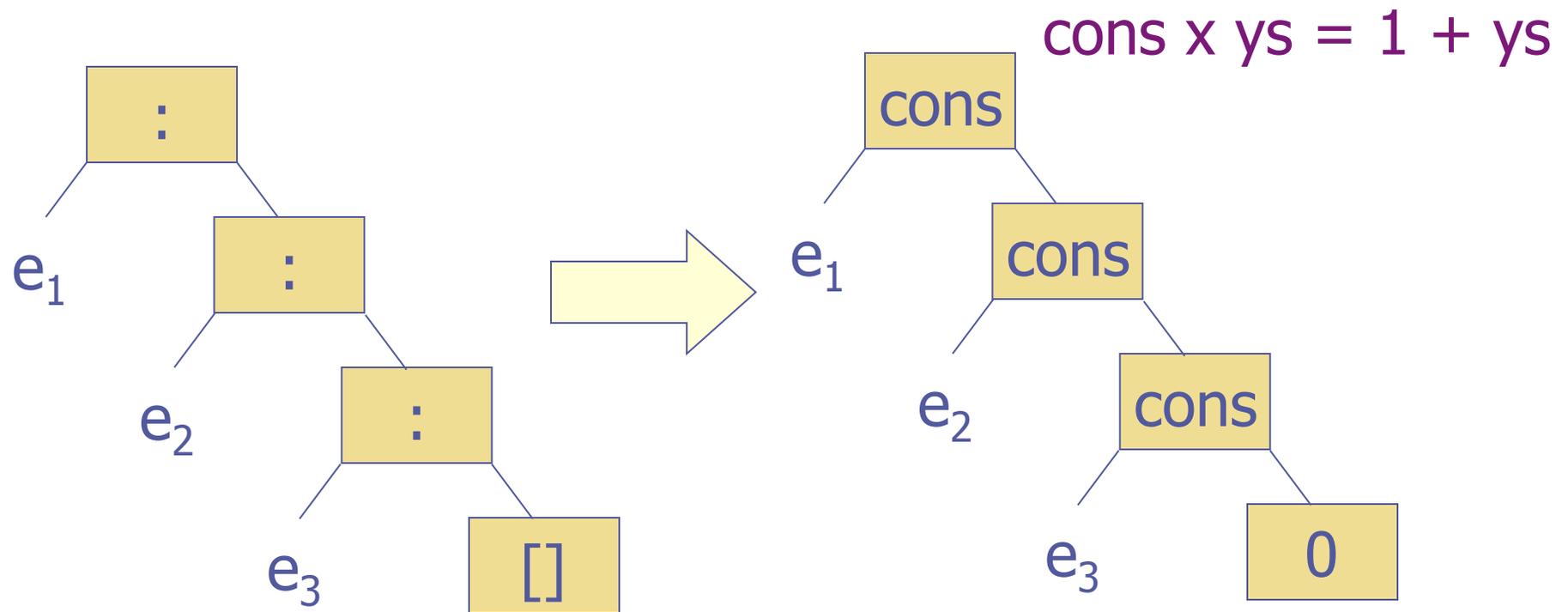
$\text{sum} = \text{foldr } (+) 0$

Example: product



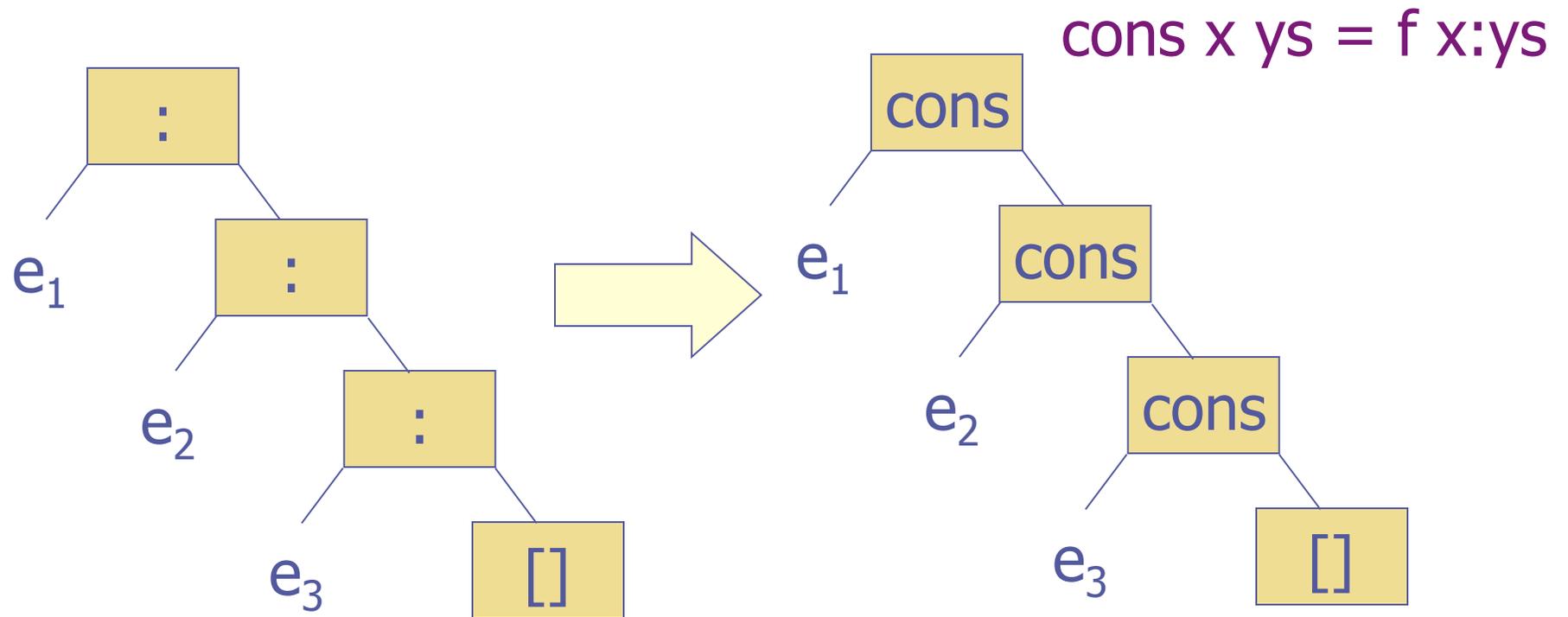
product = foldr (*) 1

Example: length



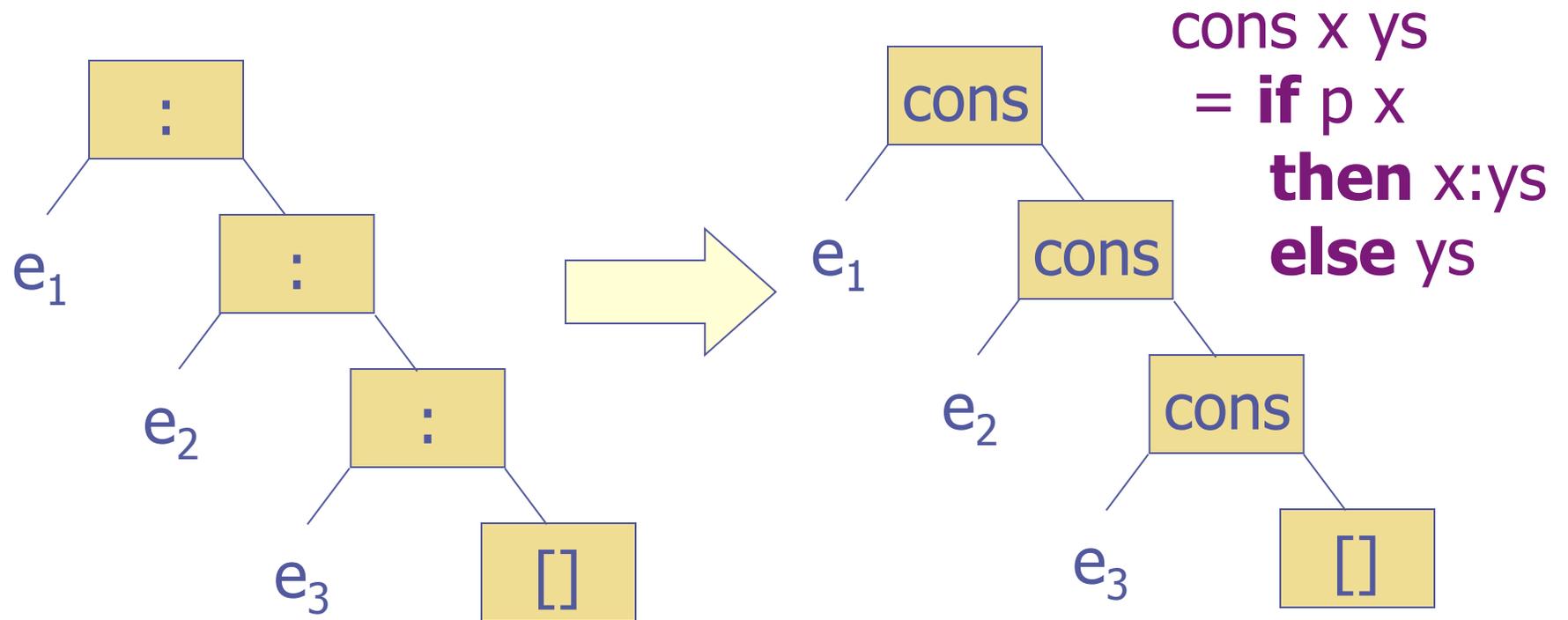
$\text{length} = \text{foldr } (\backslash x \text{ } ys \rightarrow 1 + ys) \ 0$

Example: map



$\text{map } f = \text{foldr } (\backslash x \text{ } ys \rightarrow f \ x : ys) \ []$

Example: filter



$\text{filter } p = \text{foldr } (\backslash x \text{ } ys \rightarrow \text{if } p \ x \text{ then } x:ys \text{ else } ys) \ []$

Formal Definition:

`foldr` :: $(a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b$

`foldr cons nil []` = `nil`

`foldr cons nil (x:xs)` = `cons x (foldr cons nil xs)`

Applications:

sum = foldr (+) 0

product = foldr (*) 1

length = foldr (\x ys -> 1 + ys) 0

map f = foldr (\x ys -> f x : ys) []

filter p = foldr c []

where c x ys = **if** p x **then** x:ys **else** ys

xs ++ ys = foldr (:) ys xs

concat = foldr (++) []

and = foldr (&&) True

or = foldr (||) False

Patterns of Computation:

- ◆ `foldr` captures a common pattern of computations over lists
- ◆ As such, it's a very useful function in practice to include in the Prelude
- ◆ Even from a theoretical perspective, it's very useful because it makes a deep connection between functions that might otherwise seem very different ...
- ◆ From the perspective of lawful programming, one law about `foldr` can be used to reason about many other functions

A law about foldr:

◆ If (\oplus) is an associative operator with unit n , then

$$\text{foldr } (\oplus) \ n \ xs \oplus \text{foldr } (\oplus) \ n \ ys \\ = \text{foldr } (\oplus) \ n \ (xs ++ ys)$$

◆ $(x_1 \oplus \dots \oplus x_k \oplus n) \oplus (y_1 \oplus \dots \oplus y_j \oplus n) \\ = (x_1 \oplus \dots \oplus x_k \oplus y_1 \oplus \dots \oplus y_j \oplus n)$

◆ All of the following laws are special cases:

sum xs + sum ys = sum (xs ++ ys)

product xs * product ys = product (xs ++ ys)

concat xss ++ concat yss = concat (xss ++ yss)

and xs && and ys = and (xs ++ ys)

or xs || or ys = or (xs ++ ys)

foldl:

- ◆ There is a companion function to `foldr` called `foldl`:

`foldl` $:: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b$

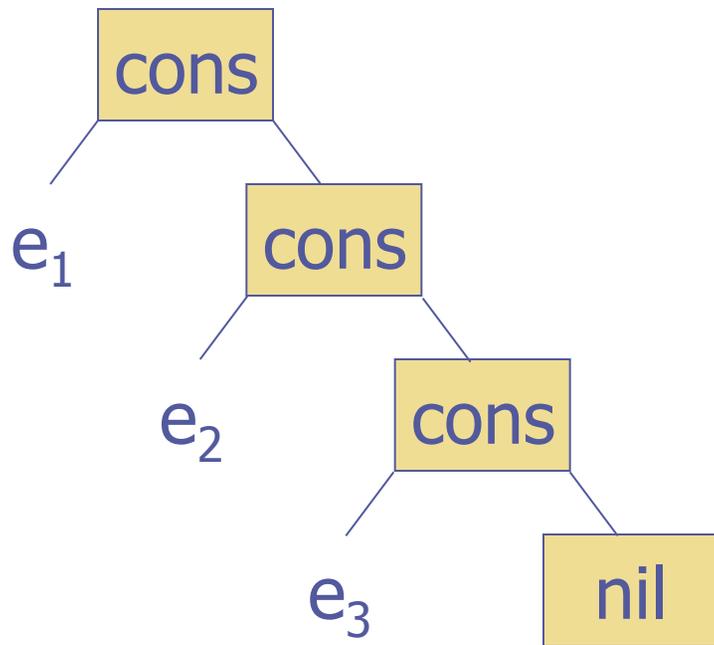
`foldl s n []` = `n`

`foldl s n (x:xs)` = `foldl s (s n x) xs`

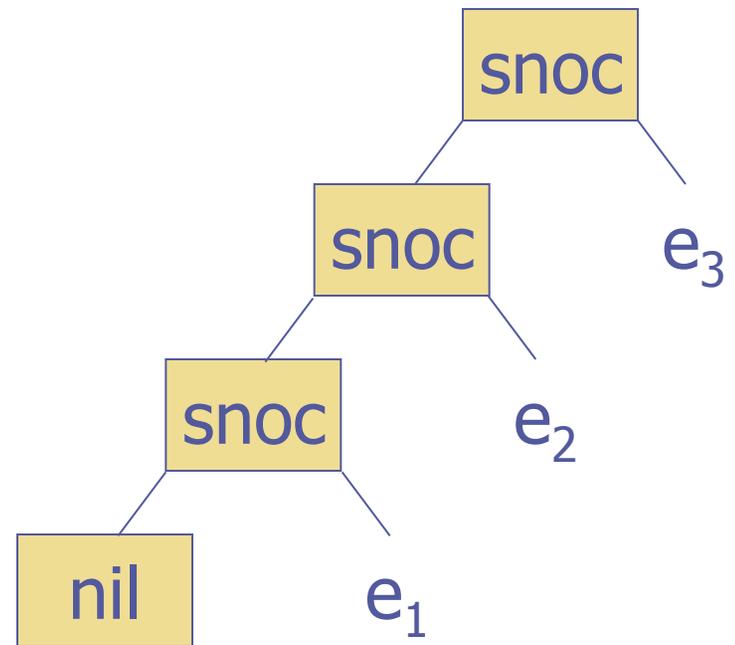
- ◆ For example:

$$\begin{aligned} & \text{foldl } s \ n \ [e_1, e_2, e_3] \\ &= s \ (s \ (s \ n \ e_1) \ e_2) \ e_3 \\ &= ((n \ `s` \ e_1) \ `s` \ e_2) \ `s` \ e_3 \end{aligned}$$

foldr vs foldl:



foldr



foldl

Uses for foldl:

- ◆ Many of the functions defined using **foldr** can be defined using **foldl**:

sum = foldl (+) 0

product = foldl (*) 1

- ◆ There are also some functions that are more easily defined using **foldl**:

reverse = foldl (\ys x -> x:ys) []

- ◆ When should you use **foldr** and when should you use **foldl**? When should you use explicit recursion instead?

foldr1 and foldl1:

- ◆ Variants of `foldr` and `foldl` that work on non-empty lists:

`foldr1` $:: (a \rightarrow a \rightarrow a) \rightarrow [a] \rightarrow a$

`foldr1 f [x]` $= x$

`foldr1 f (x:xs)` $= f\ x\ (\text{foldr1}\ f\ xs)$

`foldl1` $:: (a \rightarrow a \rightarrow a) \rightarrow [a] \rightarrow a$

`foldl1 f (x:xs)` $= \text{foldl}\ f\ x\ xs$

- ◆ Notice:

- No case for empty list
- No argument to replace empty list
- Less general type (only one type variable)

Uses of foldl1, foldr1:

From the prelude:

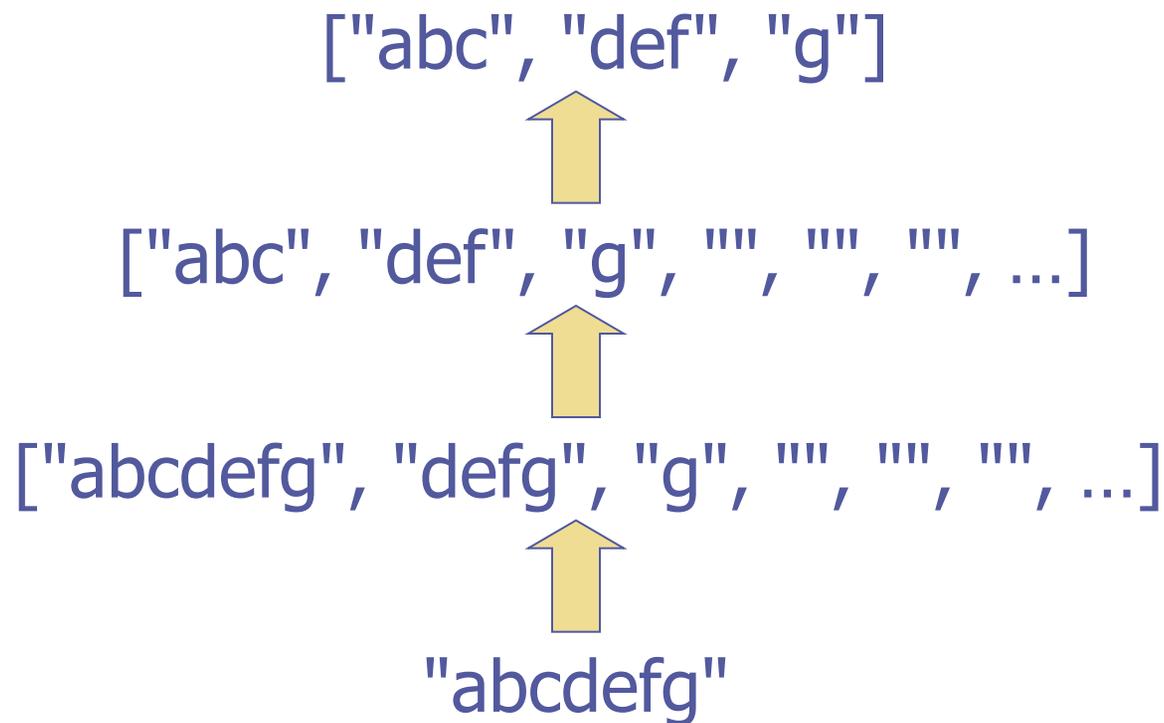
```
minimum = foldl1 min  
maximum = foldl1 max
```

Not in the prelude:

```
commaSep = foldr1 (\s t -> s ++ ", " ++ t)
```

Example: Grouping

group n = takeWhile (not.null)
· map (take n)
· iterate (drop n)



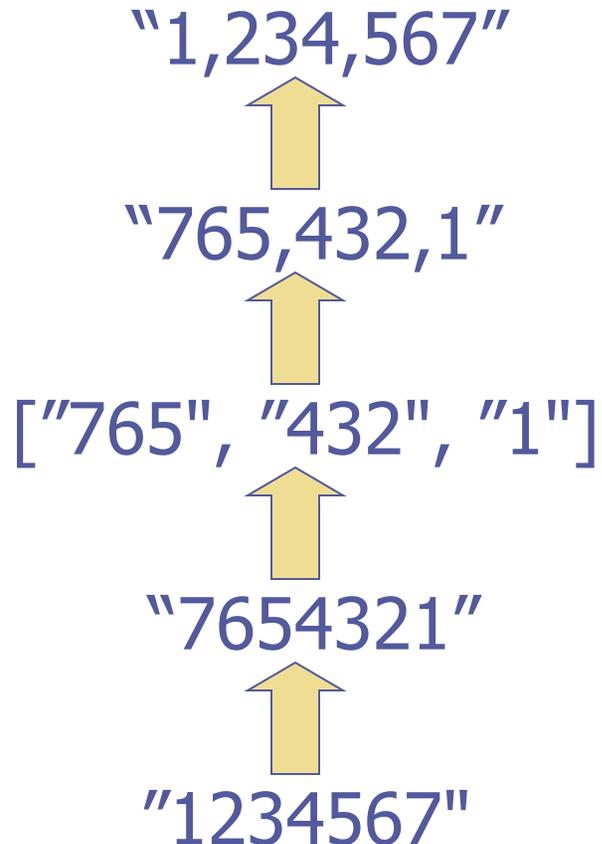
Example: Adding Commas

group n = reverse

. foldr1 (\xs ys -> xs++", "++ys)

. group 3

. reverse



Example: transpose

`transpose` :: `[[a]] -> [[a]]`

`transpose []` = `[]`

`transpose ([] : xss)` = `transpose xss`

`transpose ((x:xs) : xss)`
= `(x : [h | (h:t) <- xss])`
: `transpose (xs : [t | (h:t) <- xss])`

Example:

`transpose [[1,2,3],[4,5,6]]` = `[[1,4],[2,5],[3,6]]`

Example: say

```
Say> putStr (say "hello")
```

```
H   H   EEEEE   L   L   OOO
H   H   E       L   L   O   O
HHHHH   EEEEE   L   L   O   O
H   H   E       L   L   O   O
H   H   EEEEE   LLLLL   LLLLL   OOO
```

```
Say>
```

... continued:

```
say = ('\n':)
     . unlines
     . map (foldr1 (\xs ys->xs++" "++ys))
     . transpose
     . map picChar
```

where

```
picChar 'A' = [ "  A  ",
                 " A A ",
                 "AAAAA",
                 "A  A",
                 "A  A" ]
```

etc...

Composition and Reuse:

```
say> (putStr . concat . map say . lines . say) "A"
```

```
  A
 A A
AAAAA
 A  A
 A  A
```

```
  A          A
 A A        A A
AAAAA      AAAAA
 A  A      A  A
 A  A      A  A
```

```
  A      A      A      A      A
 A A    A A    A A    A A    A A
AAAAA  AAAAA  AAAAA  AAAAA  AAAAA
 A  A  A  A  A  A  A  A  A  A
 A  A  A  A  A  A  A  A  A  A
```

```
  A          A
 A A        A A
AAAAA      AAAAA
 A  A      A  A
 A  A      A  A
```

```
  A          A
 A A        A A
AAAAA      AAAAA
 A  A      A  A
 A  A      A  A
```

```
Say>
```

Summary:

- ◆ Folds on lists have many uses
- ◆ Folds capture a common pattern of computation on list values
- ◆ In fact, there are similar notions of fold functions on many other algebraic datatypes ...)